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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
10/658,897	09/10/2003	Vijay S. Menon	42P17027	5168
8791	7590 06/07/2006		EXAMINER	
BLAKELY SOKOLOFF TAYLOR & ZAFMAN			FLOURNOY, HORACE L	
12400 WILS SEVENTH F	HIRE BOULEVARD LOOR		ART UNIT	PAPER NUMBER
	ES, CA 90025-1030		2189	

Please find below and/or attached an Office communication concerning this application or proceeding.

	Application No.	Applicant(s)	
	10/658,897	MENON ET AL.	
Office Action Summary	Examiner	Art Unit	
	Horace L. Flournoy	2189	
The MAILING DATE of this communicate Period for Reply	ion appears on the cover sheet w	th the correspondence address	
A SHORTENED STATUTORY PERIOD FOR WHICHEVER IS LONGER, FROM THE MAIL  - Extensions of time may be available under the provisions of 37 after SIX (6) MONTHS from the mailing date of this communic.  - If NO period for reply is specified above, the maximum statutor Failure to reply within the set or extended period for reply will. I Any reply received by the Office later than three months after the earned patent term adjustment. See 37 CFR 1.704(b).	ING DATE OF THIS COMMUNION CFR 1.136(a). In no event, however, may a ration.  y period will apply and will expire SIX (6) MON by statute, cause the application to become AB	CATION.  eply be timely filed  THS from the mailing date of this communication.  ANDONED (35 U.S.C. § 133).	
Status			
<ul> <li>1) ⊠ Responsive to communication(s) filed o</li> <li>2a) ⊠ This action is FINAL. 2b)[</li> <li>3) ☐ Since this application is in condition for closed in accordance with the practice unit of the condition is in condition.</li> </ul>	☐ This action is non-final.  allowance except for formal matt	• •	
Disposition of Claims			
4) ⊠ Claim(s) <u>1-24</u> is/are pending in the appl 4a) Of the above claim(s) is/are v 5) □ Claim(s) is/are allowed. 6) ⊠ Claim(s) <u>1-24</u> is/are rejected. 7) □ Claim(s) is/are objected to. 8) □ Claim(s) are subject to restriction	vithdrawn from consideration.		
Application Papers			
9) The specification is objected to by the Entropy The drawing(s) filed on is/are: a)  Applicant may not request that any objection Replacement drawing sheet(s) including the 11) The oath or declaration is objected to by	accepted or b) objected to n to the drawing(s) be held in abeyar correction is required if the drawing	nce. See 37 CFR 1.85(a). (s) is objected to. See 37 CFR 1.121(d).	
Priority under 35 U.S.C. § 119			
12) Acknowledgment is made of a claim for a) All b) Some * c) None of:  1. Certified copies of the priority doc 2. Certified copies of the priority doc 3. Copies of the certified copies of the application from the International * See the attached detailed Office action for	cuments have been received. cuments have been received in A he priority documents have been Bureau (PCT Rule 17.2(a)).	pplication No received in this National Stage	
Attachment(s)  1)  Notice of References Cited (PTO-892)		Summary (PTO-413)	
<ul> <li>2) Notice of Draftsperson's Patent Drawing Review (PTO-3)</li> <li>Information Disclosure Statement(s) (PTO-1449 or PTO Paper No(s)/Mail Date 3/17/2006.</li> </ul>	948) Paper No(	s)/Mail Date nformal Patent Application (PTO-152)	

# **DETAILED ACTION**

### Response to Amendment

This Office action has been issued in response to amendment filed <u>17 March</u> <u>2006</u>. Claims 1-24 are pending. Applicant's arguments have been carefully and respectfully considered, but they are not entirely persuasive, as will be discussed in more detail below, even in light of the instant amendments. Accordingly, this action has been made FINAL.

### ACKNOWLEDGEMENT OF REFERENCES CITED BY APPLICANT

As required by **M.P.E.P.** 609(c), the applicant's submission of the Information Disclosure Statements dated **03/17/2006** is acknowledged by the examiner and the cited references have been considered in the examination of the claims now pending. As required by **M.P.E.P.** 609(c), a copy of the PTOL-1449 initialed and dated by the examiner is attached to the instant office action.

### Claim Rejections - 35 USC § 112

The examiner acknowledges the amendments to claims 1, 5, and 12 and the remarks on page 9 of the applicants response with regard to previous 35 USC § 112, second paragraph rejections of these claims. Accordingly the previous 35 USC § 112 rejections are herewith withdrawn.

#### **REJECTIONS BASED ON PRIOR ART**

# Claim Rejections - 35 USC § 102

The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless -

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United States.

Claims 1-24 are rejected under 35 U.S.C. 102(b) as being anticipated by <u>Kenner et al.</u>
(U.S. Patent No. 5,903,749 hereafter referred to as Kenner).

With respect to independent claims 1 and 22,

(and dependent claim 19)

"A method comprising: executing a speculative read-reordered load instruction [Kenner discloses in column 1, lines 25-26, "...these compilers must be able to freely reorder the instructions to be effective. Also see the abstract] prior to a potentially conflicting load in an instruction sequence; [disclosed, e.g. in FIG. 1, step 130. See column 11, lines 39-41] storing memory conflict information [Kenner discloses in the abstract, "storing memory conflict information"] representing the speculative ["preload array entry" (column 4, line 16) / "preload instruction" (column 2, line 37). Also see "Other Publications" of Kenner] read-reordered load; [see limitation above] executing a read-reordered load check instruction associated with the speculative read-reordered load instruction, the read-reordered load check

instruction to determine if matching an address of a the potentially conflicting load matches against an address of the stored memory conflict information; [Kenner discloses in the abstract, lines 16-18, "determining if a memory conflict has occurred between the first address and the second address using the previously stored memory conflict information."] and invalidating [Kenner discloses in column 8, lines 55-57, "FIG. 7 illustrates a circuit for selectively invalidating the memory conflict entries of the ALAT according to one embodiment of the invention."] the stored memory conflict information with a matching address [Kenner discloses in column 7, lines 61-63, "determining if a memory conflict has occurred between the first address and the second address using the previously stored memory conflict information."] to the address of the potentially conflicting load if the stored memory conflict information has a value different than a value of the potentially conflicting load." [Kenner discloses in column 17, lines 10-14, "A true conflict occurs when the load and store addresses actually overlap, while a false conflict results when the load and store addresses are actually different but the conflict resulted from the hashing scheme used." Kenner teaches a value different than a value of the potentially conflicting load.]

[With further respect to claims 22-24, Kenner discloses in <u>column 9, line 60</u>

- column 10, line 3, "...software 335 for execution..."]

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With respect to independent claim 7,

(and dependent claims 8, 17 and 19)

"A processor [See FIG. 1, element 305], comprising: a re-ordered load address table (RRLAT) to store memory conflict information representing a speculative read re-ordered load [FIG. 3, element 345: "Dynamic Memory Conflict Resolution Unit" that is executed prior to a potentially conflicting load in an instruction sequence; [See FIG. 4, step 430 and all associated text within specification]; and a monitor [Kenner discloses a "conflict resolution circuit" in column 10, paragraph 5 and FIG. 3, element 392] to: compare a potentially conflicting load against the stored memory conflict information by executing a read-reordered load check instruction associated with the speculative read-reordered load instruction, the read-reordered load check instruction to determine if an address of the potentially conflicting load matches an address of the stored memory conflict information; [Kenner discloses in column 11, lines 9-12, "At step 420, the dynamic memory conflict resolution unit 345 determines whether a memory conflict occurred involving a register associated with the check no-invalidate instruction."] and invalidate the stored memory conflict information with a matching address to the address of the potentially conflicting load if the stored memory conflict information has a value different than a value of the potentially conflicting load." [Kenner discloses this limitation, e.g. in column 3, lines 28-35, 40-45, column 17, lines 51-57]

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With respect to independent claim 16,

"A computer system, comprising: a first processor; and a second processor, [Kenner discloses in column 1, lines 19-26, "...for execution by one or more processors in the computer system...." Kenner also discloses in column 9, lines 55-58, "While this embodiment is described in relation to a single processor computer system, the invention could be implemented in a multi-processor computer system."] including: a re-ordered load address table (RRLAT) to store memory conflict information representing a speculative read reordered load received from the second processor that is executed prior to a potentially conflicting load in an instruction sequence; ["...On certain processors (e.g., superscalar processors, very long instruction word processors, etc.), allowing the compiler to reorder the object code instructions can improve performance by exposing simultaneously executable instructions. However, these compilers must be able to freely reorder the instructions to be effective." Also see "Other Publications" of Kenner] and a monitor to: compare a potentially conflicting load received from the first processor against the stored memory conflict information, by executing a read-reordered load check instruction associated with the speculative readreordered load instruction, [Kenner discloses in column 11, lines 9-12, "At step 420, the dynamic memory conflict resolution unit 345 determines whether a memory conflict occurred involving a register associated with the check no-invalidate instruction."] <u>the read-reordered load check</u> instruction to determine if an address of the potentially conflicting load matches an address of the stored memory conflict information: and [Kenner discloses Art Unit: 2189

this limitation, e.g. in column 3, lines 28-35, 40-45, column 17, lines 51-57] and to invalidate the stored memory conflict information if the stored memory conflict information has a matching address and a different value than the potentially conflicting load." [FIG. 1, step 130]

## With respect to independent claim 18,

"A computer system, comprising: a memory device; and a processor coupled to the memory device, [See FIG. 1, element 305], including: a re-ordered load address table (RRLAT) to store memory conflict information representing a speculative read re-ordered load that is executed prior to a potentially conflicting load in an instruction sequence; [See rejection of claims 1, 4, and 5]; a monitor to compare a potentially conflicting load against the stored memory conflict information by executing a read-reordered load check instruction associated with the speculative read-reordered load instruction, the read-reordered load check instruction to determine if an address of the potentially conflicting load matches an address of the stored memory conflict information: [See rejection of claim 7], and to invalidate the stored memory conflict information if the stored memory conflict information has a matching address and a different value than the potentially conflicting load [Kenner discloses in column 17, lines 42-44, "...if a matching entry is found, then the valid indication of that entry is polled. If the valid indication is in the conflict/invalid state. See also rejections to claims 2 and 3]; and a cache memory; and a memory device coupled to the processor." [Kenner discloses in column 16, lines 25-26, "This table may be Application/Control Number: 10/658,897

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implemented using any form of cache." With respect to claim 18 Kenner discloses "" in FIG. 3, element 310, "Storage Device"]

With respect to claims 2, 8, 17 and 19,

"The method of claim 1, wherein further comprising validating the stored memory conflict information with the matching address is invalidated if the stored memory conflict information has the same different value as the value of than the potentially conflicting load." [Kenner discloses in column 2, lines 55-60, "If the hardware determines that an address conflict has occurred (that the memory access required for the store instruction is the same as, or overlaps with, the memory access required by the load instruction), the check instruction will branch to correction code, which re-executes the load and any dependent instructions." Kenner further discloses this limitation in column 3, lines 28-35.]

With respect to claim 3,

"The method of claim 2, further comprising executing a wherein the validating of the stored memory conflict information of the read re-ordered load check instruction [Kenner discloses this limitation, e.g. in column 3, lines 28-35.] to determine the validity of the speculative read re-ordered load further comprises passing control by the read re-ordered load check instruction to a next instruction in the instruction sequence." [Kenner discloses this limitation e.g. in column 15, lines 14-22. Kenner further discloses in column 6, lines 7-10 "...each iteration through the loop requires an entry to be created in the

preload array in response to executing the preload instruction, as well as the invalidation of that entry in response to executing the check instruction."]

### With respect to claim 4,

"The method of claim 1, wherein the memory conflict information is stored in a read re-ordered load address table (RRLAT). " [See FIG. 3, element 390 and column 8, paragraph 1. Kenner discloses in column 8, lines 10-12, "The memory conflict resolution unit includes a conflict resolution circuit coupled to a table."]

# With respect to claims 5 and 19,

"The method of claim 4, further comprising updating the stored memory conflict information by setting a validity bit [Kenner discloses in column 3, lines 34-35, "...a valid bit indicating whether the entry currently contains valid data.], in the RRLAT [See rejection of claim 4] to a valid state when new memory conflict information is stored." [Kenner discloses in column 11, lines 55-57, "The term validity indication is used to refer to any data (e.g., a bit, a string of bits) used to identify whether memory conflict information is valid."]

#### With respect to claim 6,

"The method of claim 5, further comprising setting the validity bit to an invalid state ["invalidated"] if a later conflicting load operation [See Tables 1A and 1B] is executed [See FIG. 1]." [Kenner discloses in column 4, lines 12-17, "If the

conflict bit of this entry is set, then a memory conflict occurred and flow passes to step 140. Otherwise, flow passes to step 130...As shown in step 130, the preload array entry is invalidated by resetting its valid bit."]

With respect to claim 9,

"The processor of claim 8, wherein the <u>validating of the</u> stored memory conflict information of the read re-ordered load check instruction further includes is invalidated if the stored memory conflict has a matching address and a different value than the potentially conflicting load passing control by the read re-ordered load check instruction to a next instruction in the instruction sequence. [Kenner discloses this limitation e.g. in <u>column 15</u>, lines 14-22. Kenner discloses in <u>column 17</u>, lines 42-44, "...if a matching entry is found, then the valid indication of that entry is polled. If the valid indication is in the conflict/invalid state. See <u>also column 19</u>, lines 5-29 and the rejection to claims 2 and 8]

With respect to claim 10,

"The processor of claim 7, wherein the RRLAT [See rejection of claims 1 and 4] is referenced [Note: RRLAT is part of element 345 of FIG. 3. See also the rejection to claims 4 and 5] upon the execution of a read re-ordered load check instruction to determine the validity of the speculative read re-ordered load." [See rejection of claim 20 below]

With respect to claim 11,

"The processor of claim 7, wherein the RRLAT [See rejection of claims 1 and

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4] may be any one of a direct-mapped, multi-way set associative [Kenner

discloses in column 16, line 54, "set-associative"], and fully associative data

structure [See column 16, paragraphs 2-4]." [Kenner further discloses in

column 16, lines 25-26, "This table may be implemented using any form of

cache."]

With respect to claim 12,

"The processor of claim 7, wherein the RRLAT [FIG. 7 and also rejection of

claims 1 and 4] is portioned among hardware thread contexts." [Kenner

discloses in column 18, lines 3-7, "The dynamic memory conflict resolution

unit 705 is shown including a table 710 for storing the memory conflict

entries. Although only one storage area in the table 710 is being described,

each of the storage areas in the table 710 has associated with it a set of

comparators (740) and an AND gate (745). " The examiner interprets

"hardware thread contexts" as analogous to comparators and an AND

gate.]

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With respect to claim 13,

"The processor of claim 7, wherein the RRLAT [See rejection of claims 1 and 4] includes storage locations for an address [Kenner discloses in column 18, lines 7-10, "Each memory conflict entry stored in table 710...], a target register ID ["...has a location identifier field (715),], a value ["...a signature field (720), a width field (722),], and validity information associated with the speculative read re-ordered load." ["...and a valid indication field (725). See FIG. 7, elements 710, 715, 720, 722, and 725]

With respect to claims 20 and 24,

"The machine-readable medium of claim 23, wherein the validating of the stored memory conflict information of the read re-ordered load check instruction further comprises passing control by the read re-ordered load check instruction to a next instruction in the instruction sequence the sequence of instructions when executed by the computer system further causing the computer system to validate stored memory conflict information with a matching address if the stored memory conflict information has a matching value to the potentially conflicting lead." [Kenner discloses this limitation e.g. in column 15, lines 14-22]

With respect to claim 21,

"The computer system of claim 18, further comprising a bus [See FIG. 3, element 315, "Bus"] to control communications [Kenner discloses in column 9, lines 53-55, "The bus 315 represents one or more busses (e.g., PCI, ISA, X-Bus, EISA, VESA, etc.) and bridges (also termed as bus controllers)."]

between the processor [See FIG. 3, element 305, "Processor"] and the memory device." [See FIG. 3, element 310, "Storage Device"]

With respect to claim 23,

"The machine-readable medium of claim 22, the sequence of instructions [Kenner discloses in the abstract, "sequences of instructions"], when executed by the computer system ["When executed by a computer system, the sequences of instructions cause the computer system to perform a series of steps..."], further causing the computer system to validate stored memory conflict information with a matching address if the stored memory conflict information has the same a matching value as the value of [[to]] the potentially conflicting load." [Kenner discloses in column 2, lines 55-60, "If the hardware determines that an address conflict has occurred (that the memory access required for the store instruction is the same as, or overlaps with, the memory access required by the load instruction), the check instruction will branch to correction code, which re-executes the load and any dependent instructions."]

[Note: With respect to claims 22-24, Kenner discloses a machinereadable medium in the abstract.] Application/Control Number: 10/658,897

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### **ACKNOWLEDGMENT OF ISSUES RAISED BY THE APPLICANT**

# Response to Amendment

Applicant's arguments filed **March 17, 2006** have been fully considered but they are not deemed to be persuasive and, as required by **M.P.E.P. 707.07(f)**, a response to these arguments appears below.

#### **ARGUMENTS CONCERNING PRIOR ART REJECTIONS**

# 1<sup>ST</sup> POINT OF ARGUMENT:

With respect to the arguments on page 10 of the applicant's remarks, the examiner respectfully disagrees with the applicant that Kenner does not disclose or suggest <u>invalidating the stored memory conflict information with a matching address to the address of the potentially conflicting load if the stored memory conflict information has a value different than a value of the potentially conflicting load. Kenner discloses in column 17, lines 10-14, "A true conflict occurs when the load and store addresses actually overlap, while a false conflict results when the load and store addresses are actually different but the conflict resulted from the hashing scheme used." See also: column 3, lines 28-35, 40-45, column 17, lines 51-57. The examiner finds that Kenner teaches the limitation above. See rejection of claim 1 supra.</u>

The examiner notes that <u>column 3</u>, <u>lines 28-35</u> lists information, which is all interpreted as **a value** for the stored memory conflict information. See FIG. 4 and all associated text within the specification.

Furthermore, with regard to the arguments on page 11 of the applicant's response, independent claims 7, 16, 18, and 22 are all rejected on the same ground as independent claim 1. Kenner anticipates the limitations of these claims, as stated in the above rejections.

#### CONCLUSION

THIS ACTION IS MADE FINAL. Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the mailing date of this final action.

# **Direction of Future Correspondences**

Any inquiry concerning this communication or earlier communication from the examiner should be directed to Horace L. Flournoy whose telephone number is (571)

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272-2705. The examiner can normally be reached on Monday through Friday 8:00 AM

to 5:30 PM (ET).

**Important Note** 

If attempts to reach the examiner by telephone are unsuccessful, the examiner's

supervisor, Reginald G. Bragdon can be reached on (571) 272-4204. The fax phone

numbers for the organization where this application or proceeding is assigned is (703)

746-7239.

Information regarding the status of an Application may be obtained from the

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questions on access to the Private PAIR system, contact the Electronic Business

Center (EBC) at 866-217-9197 (toll-free).

Any inquiry of a general nature or relating to the status of this application or

proceeding should be directed to the receptionist whose telephone number is (571) 272-

2100.

Horace L. Flournoy

teynald B. Bragdon
Reginald G. Bragdon

Patent Examiner Art unit: 2189

Supervisory Patent Examiner
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